

# Finding All Sorting Tandem Duplication Random Loss Operations

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# Outline

- 1 Introduction
- 2 All Sorting TDRLs
  - A Restricted Case
  - The General Case
- 3 Late Breaking Findings
- 4 Results

# Genome Rearrangement

>species1

.. G1 G2 G3 G4 G5 G6 G7 ..

>species2

.. G4 G2 G7 G1 -G3 G5 G6 ..

⋮

>species $m$

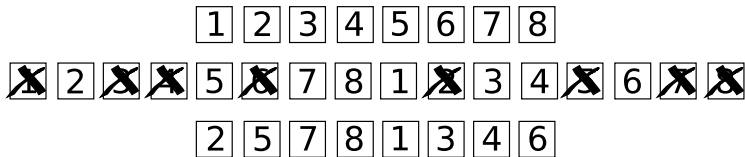
.. G7 G1 G2 G6 G5 G4 G3 ..



[www.tolweb.org](http://www.tolweb.org)

Gene arrangements = permutations.

# Tandem Duplication Random Loss (TDRL)

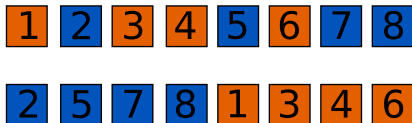


## Definition (TDRL)

TDRL  $\tau(F, S)$  defined by:

- $F$  the set of elements kept in the first copy and
- $S$  the set of elements kept in the second copy.

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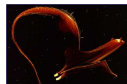
## Relevance of TDRs

*S. Purpuratus*

ND4L CO2 K ATP8 ATP6 CO3 -S2 ND3 ND4 H S1 ND5 -ND6 CYTB F 12S E T P -Q N L1 -A W G -Y  
E P N L1 W -Y ND4L CO2 K ATP8 ATP6 CO3 -S2 ND3 ND4 H S1 ND5 -ND6 CYTB F 12S T -Q -A C

*C. miniata**C. Sloani*

ATP8 ATP6 CO3 G ND3 R ND4L ND4 H S1 L1 ND5 -ND6 -E CYTB T -P E 12S V 16S L2 ND1 I -Q M ND2 W  
R ND4L ND4 H S1 -E CYTB F 12S V 16S L2 ND1 -Q W ATP8 ATP6 CO3 G ND3 L1 ND5 -ND6 T -P I M ND2

*E. Plecanoides**S. fontinalis*

-Q M ND2 V -A -N -C S1 CO1 S2 D CO2 K ATP8 ATP6 CO3 G ND3 R ND4L ND4 H S1 S1 ND5 ND6 -E CYTB T -P F 12S V 16S S2 ND1  
ND2 W -Y CO1 S2 D CO2 K ATP8 ATP6 CO3 G ND3 R ND4L ND4 H L1 ND5 ND6 P F 12S V 16S ND1 -Q M -A -N -C S1 -E CYTB T L2

*B. nectabanus**L. polyphemus*

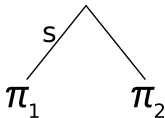
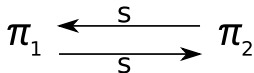
ND3 A R N S2 E -F -ND5 -H -ND4 -ND4L T -P ND6 COB S -ND1 -L2 -L -RNL -V -RNS I -Q M ND2 W -G -Y  
A R S2 E -F -ND5 -ND4L ND6 S -ND1 M -C -Y ND3 N -H -ND4 T -P COB -L2 -L -RNL -V -RNS I -Q ND2 W

*S. coleoprata*

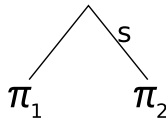
# Relevance of TDRLs

Asymmetry provides additional phylogenetic information

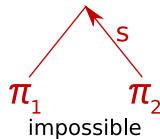
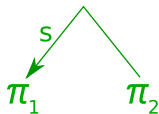
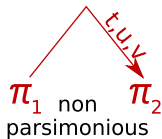
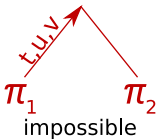
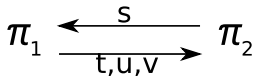
Symmetric operation:  
(Inversion, Transposition, ...)



?



Asymmetric operation:  
(TDRL)



# Sorting by TDRLs

## Definition (Sorting by TDRLs)

Given a permutation  $\pi$ . Find a shortest sequence of TDRLs

$\tau_1, \dots, \tau_{d(\pi)}$  such that  $\pi \circ \tau_1 \circ \dots \circ \tau_{d(\pi)} = \mathbf{1}$ .

The **TDRL distance** is the length of the sequence, denoted by  $d(\pi)$ .

## Definition (Chain of a permutation $\pi$ )

A chain of a permutation  $\pi$  is a maximal list  $(e_1, \dots, e_k)$  of elements of  $\pi$  where

$e_{i+1} = e_i + 1$  and  $\pi^{-1}(e_i) < \pi^{-1}(e_{i+1})$ .

Number of chains of  $\pi$ :  $\rho(\pi)$ .

Sorting  $\equiv$  Merge the chains in order to get one chain.

Indexing Scheme:  $c < c'$  iff  $\forall e \in c, \forall e' \in c' : e < e'$



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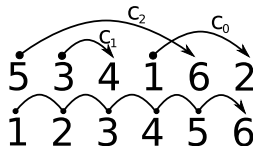
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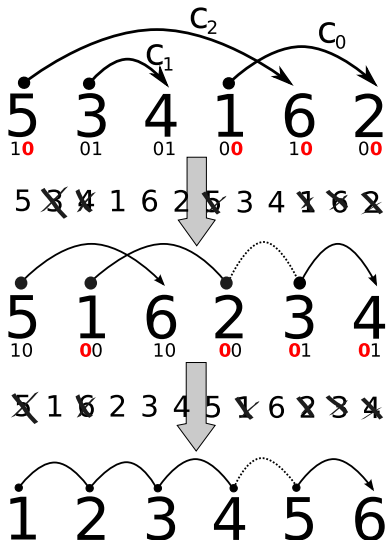
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# Sorting by TDRLs [Chaudhuri et al., SODA, 2006]



TDRL Distance:

$$d(\pi) = \lceil \log_2(\rho(\pi)) \rceil$$

Radix-Sort inspired algorithm:

- Get the binary representation of the chain index of each element
- In the  $i$ -th step: keep the elements of chains with a 0 at the  $i$ -th least significant bit in the first copy

# All Sorting TDRs

## Problem Definition

**Question 1:** Are there alternative sorting TDRL scenarios?

**Definition (All Sorting TDRs)**

Find the set of TDRs  $\{\tau : d(\pi \circ \tau) < d(\pi)\}$ .

**Question 2:** How many sorting TDRs are there?

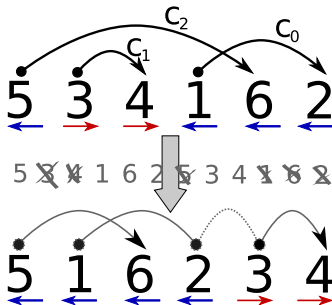
**Definition (Number of Sorting TDRs)**

Determine  $|\{\tau : d(\pi \circ \tau) < d(\pi)\}|$ .

# Basic Properties

## Observations

- *Elements kept in the 1st (2nd) copy are moved to the left (right)*
- *The order of the elements kept in the same copy is not changed*



# Restricted TDRLs

## Definition (Restricted TDRL)

All elements of a chain are kept in the same copy.

## Proposition

*Two chains  $c_i$  and  $c_j$  can be connected with a TDRL iff  $j = i + 1$ .*

*This can be done by keeping the elements of  $c_i$  in the 1st copy and the elements of  $c_{i+1}$  in the 2nd copy.*

## Proposition

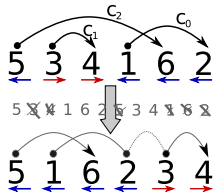
*Three chains  $c_i$ ,  $c_{i+1}$ , and  $c_{i+2}$  can not be connected at once.*

$\Rightarrow$  *Restricted TDRL distance = TDRL distance*



# Reformulation

- Restricted TDRL  $\equiv$   
 Binary string  $s$  of length  $\rho(\pi)$ :  
 $s_i = 1 \leftrightarrow c_i$  is kept in the 1st copy  
 $s_i = 2 \leftrightarrow c_i$  is kept in the 2nd copy
- $c_i$  and  $c_{i+1}$  get connected iff  $s_i s_{i+1} = 12$



$$\frac{s_0 \quad s_1 \quad s_2}{1 \quad 2 \quad 1}$$

Question:

- How many strings of length  $n$  with at least  $k$  12-transitions?

# String Count

Number of strings of length  $n$  which have exactly  $k$  12-transitions.

$$\begin{array}{ll}
 s_1 = 1 \wedge s_n = 1 & \rightarrow \binom{n-1}{2k} & \begin{array}{cccc} 1 & \blacktriangleup & 2 & \blacktriangledown & 1 & 1 & \blacktriangleup & 2 & 2 & \blacktriangledown & 1 & 1 \end{array} \\
 s_1 = 1 \wedge s_n = 2 & \rightarrow \binom{n-1}{2k-1} & \begin{array}{cccc} 1 & \blacktriangleup & 2 & \blacktriangledown & 1 & 1 & \blacktriangleup & 2 & 2 & \blacktriangledown & 1 & \blacktriangleup & 2 \end{array} \\
 s_1 = 2 \wedge s_n = 1 & \rightarrow \binom{n-1}{2k+1} & \begin{array}{cccc} 2 & 2 & \blacktriangledown & 1 & 1 & \blacktriangleup & 2 & 2 & \blacktriangledown & 1 & 1 \end{array} \\
 s_1 = 2 \wedge s_n = 2 & \rightarrow \binom{n-1}{2k} & \begin{array}{cccc} 2 & 2 & \blacktriangledown & 1 & 1 & \blacktriangleup & 2 & 2 & \blacktriangledown & 1 & \blacktriangleup & 2 \end{array}
 \end{array}$$

$$= \binom{n+1}{2k+1}$$

Number of strings of length  $n$  which have at least  $k$  12-transitions.

$$= \sum_{i=k}^{\lfloor \frac{n}{2} \rfloor} \binom{n+1}{2i+1}$$

# Result

## Theorem

For a permutation  $\pi$  with  $\rho$  chains there are

$$\sum_{i=\rho-2^{\lceil \log_2(\rho) \rceil - 1}}^{\lfloor \frac{\rho}{2} \rfloor} \binom{\rho+1}{2i+1}$$

sorting restricted TDRLs.

For  $\rho = 2^x$ :

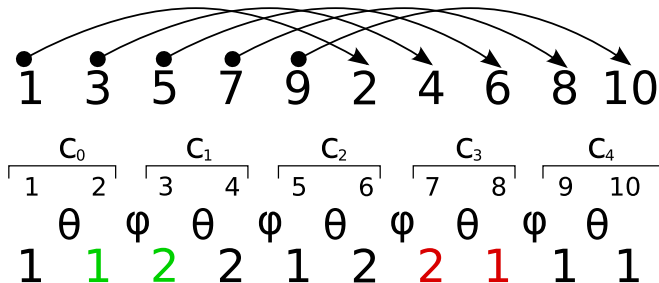
- Only one sorting TDRL
- Only one sorting TDRL scenario

In general:

- Each sorting scenario is unique after  $\lceil \log_2(2^{\lceil \log_2(\rho) \rceil} - \rho + 1) \rceil$  sorting TDRLs.



## All Sorting TDRLs

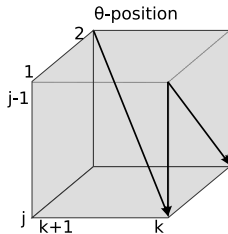
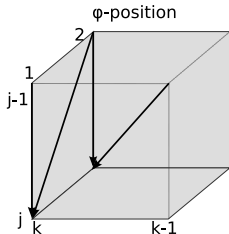


- 21 at  $\theta$ -positions breaks a chain  $\rightarrow \rho$  increased by 1
- 12 at  $\phi$ -positions connect chains  $\rightarrow \rho$  decreased by 1

**Question:** How many binary strings of length  $n$  with  $k \leq \Delta\rho$ .

# Dynamic Programming Approach

$a_{j,k}^x$ : Number of possible binary strings of length  $j$  ending with  $x \in \{1, 2\}$  that change the number of chains by  $k$ .



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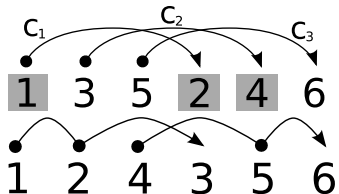
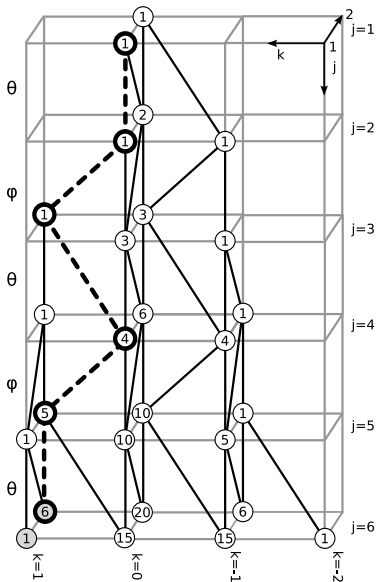
$$a_{j,k}^2 = \begin{cases} a_{j-1,k+1}^1 + a_{j-1,k}^2 & \text{if } p_{j-1} = \phi \\ a_{j-1,k}^1 + a_{j-1,k}^2 & \text{else} \end{cases}$$

$$a_{j,k}^1 = \begin{cases} a_{j-1,k}^1 + a_{j-1,k-1}^2 & \text{if } p_{j-1} = \theta \\ a_{j-1,k}^1 + a_{j-1,k}^2 & \text{else} \end{cases}$$

$a_{1,0}^1 = 1$ ,  $a_{1,0}^2 = 1$  other values initialised to 0

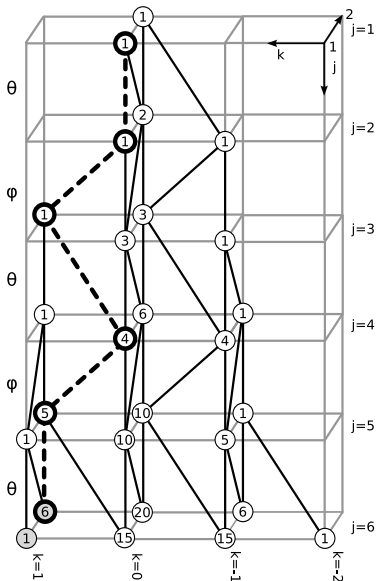
Number of sorting TDRL:  $\sum_{i=\rho-2^{\lceil \log_2(\rho) \rceil - 1}}^{\lfloor \frac{\rho}{2} \rfloor} a_{n,i}^1 + a_{n,i}^2$

# Dynamic Programming Approach



Sorting TDRL events can be enumerated by backtracking.

# Equalities



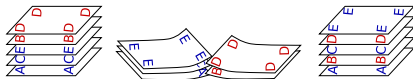
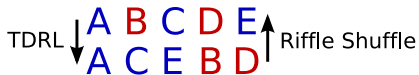
Number of sorting TDRs:

$$\sum_{i=0}^{2^{\lceil \log_2(\rho(\pi)) \rceil} - \rho(\pi)} \binom{n}{i}$$

Number of sorting restricted TDRs:

$$\sum_{i=0}^{2^{\lceil \log_2(\rho(\pi)) \rceil} - \rho(\pi)} \binom{\rho(\pi)}{i}$$

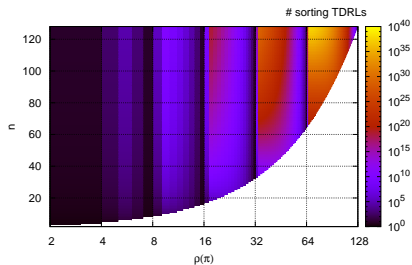
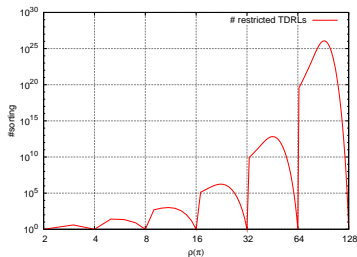
# Riffle Shuffle



Riffle Shuffle Distance:

- Schwenk *Elementary Problem: E3143*, Am. Math Mon., 1986
- Schwenk *E3143*, Am. Math Mon., 1988

# Number of Sorting TDRLs

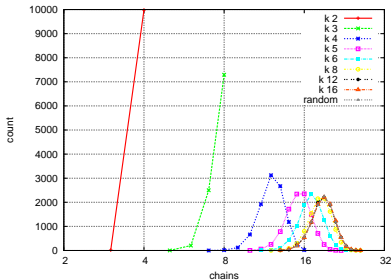


- Only 1 sorting TDRL for  $\rho = 2^x$

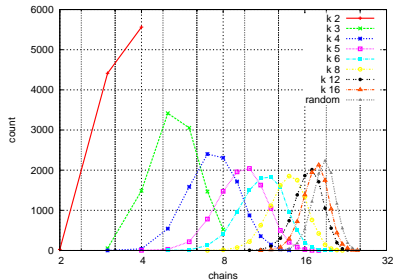
# Number of Chains in Simulated Permutations

- Apply  $k$  “random” TDRs on  $\tau$

### Random TDRs



### Random TDRs in random interval



- Unique sorting scenarios for  $\rho \in \{2, 4\}$  (and  $\rho = 8$ ) are very likely



# Mitochondrial Data

Previously unpublished scenario of two TDRLs:



*S. fontinalis*



Support:

- 4 chains → unique scenario
- *P. myriaster* → *S. fontinalis* needs 3 TDRLs
- Reversal distance is 15
- Transposition distance is 7
- There exists no TDRL median with score 2 or less
- Fragments of duplicated sequences support TDRLs [Miya05]

*P. myriaster*



# Conclusion

- Method for **enumerating all sorting TDRLs**
- **Closed formulas** for the number of sorting TDRLs
- Identification of **unique** TDRL sorting scenarios possible
- Identification of likely unique TDRLs scenario in mitochondria

Thank You!

